

# Knowledge Representation and Reasoning Logics for Artificial Intelligence

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# 12 Production Systems Architecture

## Working (Short-term) Memory

Contains set (unordered, no repeats) of  
Working Memory Elements (WMEs).  
Each being a rather flat, ground (no variables) symbol structure.

# Rule (Long-term) Memory

Contains set (unordered, no repeats) of  
Production Rules.

Each being a condition-action rule  
of form

**if** condition<sub>1</sub> ... condition<sub>*n*</sub> **then** action<sub>1</sub> ... action<sub>*m*</sub>

Each condition and action being like a WME,  
but allowing variables (and, maybe, other expressions)

# Rule Triggering

A rule **if** condition<sub>1</sub> ... condition<sub>n</sub> **then** action<sub>1</sub> ... action<sub>m</sub>  
is triggered  
if there is a substitution,  $\sigma$   
such that each condition<sub>i</sub> $\sigma$  is a WME.

A single rule can be triggered in multiple ways (by multiple substitutions).

# Rule Firing

A rule **if** condition<sub>1</sub> ... condition<sub>n</sub> **then** action<sub>1</sub> ... action<sub>m</sub>  
that is triggered in a substitution  $\sigma$   
fires by performing every action<sub>i</sub> $\sigma$ .

# Production System Execution Cycle

**loop**

Collect  $\mathcal{T} = \{r\sigma \mid r\sigma \text{ is a triggered rule}\}$

**if**  $\mathcal{T}$  is not empty

    Choose a  $r\sigma \in \mathcal{T}$

    Fire  $r\sigma$

**until**  $\mathcal{T}$  is empty.



# Some Typical Actions

- stop
- delete a WME
- add a WME
- modify a WME
- formatted print

# Conflict Resolution Strategies

Purpose: to “Choose a  $r\sigma \in \mathcal{T}$ ”

*Specificity*: If the conditions of one rule are a subset of a second rule, choose the second rule. [B & L, p. 126]

*Recency*: Based on recency of addition or modification of WMEs, or on recency of a rule firing. [B & L, p. 126]

*Refactoriness*: Don't allow the same substitution instance of a rule to fire again. [B & L, p. 127]

*Saliency*: Explicit saliency value. “The use of saliency is generally discouraged” [<http://herzberg.ca.sandia.gov/jess/docs/70/rules.html#saliency>].

# The Rete Algorithm

## Assumptions

Rule memory doesn't change.

WM changes only slightly on each cycle.

WMEs are ground.

Production Systems are data-driven (use forward chaining).

Many rules share conditions.

# The Rete Network

Create a network from the conditions (Like a discrimination tree) with rules at the leaves.

Create a token for each WME.

Pass each token through the network, stopping when it doesn't satisfy a test; resuming when the WME is modified.

When tokens reach a leaf, the rule is triggered.

Kinds of branch nodes

*$\alpha$  nodes:* Simple test.

*$\beta$  nodes:* Constraints caused by different conditions.