

CSE 431/531: Algorithm Analysis and Design (Spring 2020)

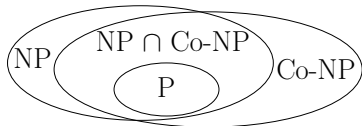
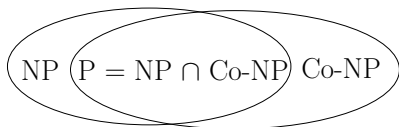
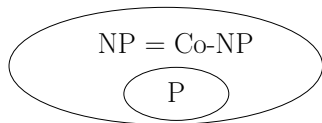
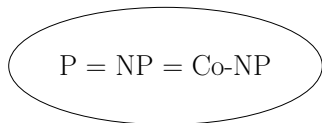
## NP-Completeness – Recitation

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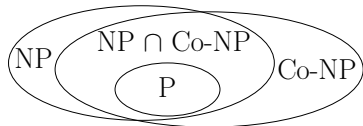
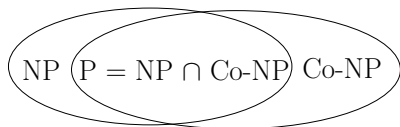
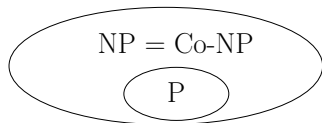
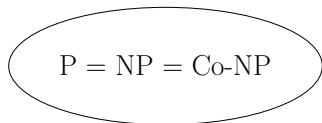
# Exercises

Recall the 4 scenarios:



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- Prove:  $P = NP$  if and only if  $P = \text{Co-NP}$

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For each of the following problem  $X$ , answer: whether (1)  $X \in \text{NP}$ , (2)  $X \in \text{CO-NP}$ . Each answer is either “yes” or “we do not know”.

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- 2 Given a graph  $G = (V, E)$  and an integer  $t > 0$ , whether the **minimum** vertex cover of  $G$  has size **at least**  $t$ .

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- 3 Given a directed graph  $G = (V, E)$ , with weights  $w : E \rightarrow \mathbb{R}_{>0}$ ,  $s, t \in V$ , and a number  $L > 0$ , whether the length of the shortest path from  $s$  to  $t$  in  $G$  is at most  $L$ .

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- 4 Given two boolean formulas, whether they are equivalent.  
For example,  $(x_1 \vee x_2) \wedge (\neg x_1 \vee x_3)$  and  $(\neg x_1 \wedge x_2) \vee (x_1 \wedge x_3)$  are equivalent since they give the same value for every assignment of  $(x_1, x_2, x_3)$ .

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- 3 Given a graph  $G = (V, E)$ , the degree-3 spanning tree (D3ST) problem asks whether  $G$  contains a spanning tree  $T$  of degree at most 3. Prove Hamiltonian-Path  $\leq_P$  D3ST.