CSE 431/531: Algorithm Analysis and Design (Fall 2022)
Divide-and-Conquer – Recitation

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For each of the following recurrences, use the master theorem to give the tight asymptotic upper bound.

1. \( T(n) = 4T(n/3) + O(n) \).
   \( T(n) = O(\ ) \)

2. \( T(n) = 3T(n/3) + O(n) \).
   \( T(n) = O(\ ) \)

3. \( T(n) = 4T(n/2) + O(n^2 \sqrt{n}) \).
   \( T(n) = O(\ ) \)

4. \( T(n) = 8T(n/2) + O(n^3) \).
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Solving Recurrences

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   \( T(n) = O(n^3 \log n) \)
Covering Chessboard using L-shape Tiles

Consider a $2^n \times 2^n$ chessboard with one arbitrary chosen square removed. Prove that any such chessboard can be tiled without gaps by L-shaped pieces, each composed of 3 squares. The following figure shows how to tile a $4 \times 4$ chessboard with the square on the left-top corner removed, using 5 L-shaped pieces.
Finding Local Minimum In a 1-D Array

Given an array $A[1 .. n]$ of $n$ distinct numbers, we say that some index $i \in \{1, 2, 3 \cdots , n\}$ is a local minimum of $A$, if $A[i] < A[i - 1]$ and $A[i] < A[i + 1]$ (we assume that $A[0] = A[n + 1] = \infty$).

Suppose the array $A$ is already stored in memory. Give an $O(\log n)$-time algorithm to find a local minimum of $A$. 

Given two $n$-digit integers, output their product. Design an $O(n^{\log_2 3})$-time algorithm to solve the problem. Notice that you cannot multiply two big integers directly using a single operation.
Majority and Weak Majority

Given an array of integers $A[1..n]$, we would like to decide if

1. there exists an integer $x$ which occurs in $A$ more than $n/2$ times. Give an algorithm which runs in time $O(n)$.
2. there exists an integer $x$ which occurs in $A$ more than $n/3$ times. Give an algorithm which runs in time $O(n)$.

You can assume we have the algorithm Select as a black-box, which, given an $n$-size array $A$ and integer $1 \leq i \leq n$, can return the $i$-th smallest element in a size $n$-array in $O(n)$-time.
Given two sorted arrays $A$ and $B$ with total size $n$, you need to design and analyze an $O(\log n)$-time algorithm that outputs the median of the $n$ numbers in $A$ and $B$. You can assume $n$ is odd and all the numbers are distinct. For example,

- **Input:** $A = [3, 5, 12, 18, 50],$
- $B = [2, 7, 11, 30],$
- **Output:** 11
- **Explanation:** the merged set is $[2, 3, 5, 7, 11, 12, 18, 30, 50]$