CSE 486/586 Distributed Systems Consistency --- 2

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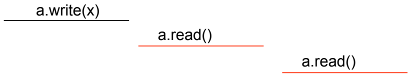
Recap: Linearizability

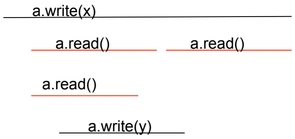
- Linearizability
  - Should provide the behavior of a single copy
  - A read operation returns the most recent write, regardless of the clients.
  - “The most recent”: determined by time.
- Complication
  - In the presence of concurrency, read/write operations overlap.

Recap: Linearizability Complications

- Non-overlapping ops: time-based clear-cut ordering

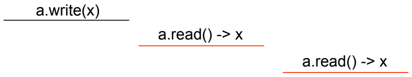


- Overlapping ops: not clear-cut with time

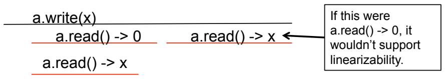


Linearizability Examples

- Example 1

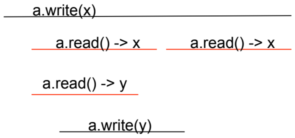


- Example 2


a.write(x)
a.read() -> 0
a.read() -> x
a.read() -> x
If this were a.read() -> 0, it wouldn’t support linearizability.

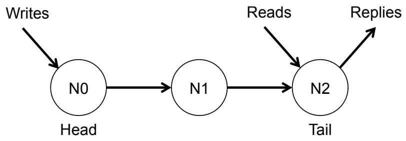
Linearizability Examples

- Example 3

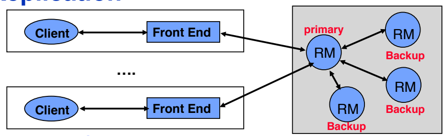


Chain Replication

- One technique to provide linearizability



Passive (Primary-Backup) Replication



- Request Communication: the request is issued to the primary RM and carries a unique request id.
- Coordination: Primary takes requests atomically, in order, checks id (resends response if not new id.)
- Execution: Primary executes & stores the response
- Agreement: If update, primary sends updated state/ result, req-id and response to all backup RMs (1-phase commit enough).
- Response: primary sends result to the front end

CSE 486/586 Administrivia

- PA3 deadline: 4/11 (Friday)

Linearizability vs. Sequential Consistency

- Both care about giving an illusion of a single copy.
  - From the outside observer, the system should (almost) behave as if there's only a single copy.
- Linearizability cares about time.
  - Steve writes on his facebook wall at 11am.
  - Atri writes on his facebook wall at 11:05am.
  - Everyone will see the posts in that order.
- Sequential consistency cares about program order.
  - Steve writes on his facebook wall at 11am.
  - Atri writes on his facebook wall at 11:05am.
  - It's not necessarily that the posts will be ordered that way (though everyone will see the same order).

Sequential Consistency

- Sequential consistency
  - Should provide the behavior of a single copy
  - A read operation returns the most recent write, regardless of the clients.
- "most recent"
  - Ops within the same client: determined by time (program order)
  - Ops across clients: Not determined by time, i.e., we can re-order them.
  - I.e., we just need to preserve the program order

Sequential Consistency

- To the outside observer, the system needs to provide a global ordering of operations where:
  - It works like a single copy.
  - The ordering of ops coming from the same client is preserved.
- Linearizability vs. sequential consistency
  - With sequential consistency, the system has freedom as to how to interleave operations coming from different clients, as long as the ordering from each client is preserved.
  - With linearizability, the interleaving across all clients is pretty much determined already based on time.

Sequential Consistency Examples

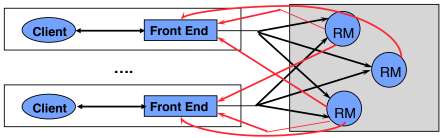
- Example 1

```
– P1: a.write(A)
– P2:             a.write(B)
– P3:                        a.read()->B     a.read()->A
– P4:                               a.read()->B     a.read()->A
```

- Example 2

```
– P1: a.write(A)
– P2:             a.write(B)
– P3:                        a.read()->B     a.read()->A
– P4:                               a.read()->A     a.read()->B
```

Active Replication



- Request Communication: The request contains a unique identifier and is multicast to all by a reliable totally-ordered multicast.
- Coordination: Group communication ensures that requests are delivered to each RM in the same order (but may be at different physical times!).
- Execution: Each replica executes the request. (Correct replicas return same result since they are running the same program, i.e., they are replicated protocols or replicated state machines)
- Agreement: No agreement phase is needed, because of multicast delivery semantics of requests
- Response: Each replica sends response directly to FE

Two More Consistency Models

- Even more relaxed
  - We don't even care about providing an illusion of a single copy.
- Causal consistency
  - We care about ordering causally related write operations correctly.
- Eventual consistency (next lecture)
  - As long as we can say all replicas converge to the same copy eventually, we're fine.

Summary

- Linearizability
  - The ordering of operations is determined by time.
  - Primary-backup can provide linearizability.
  - Chain replication can also provide linearizability.
- Sequential consistency
  - The ordering of operations preserves the program order of each client.
  - Active replication can provide sequential consistency.

Acknowledgements

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