

## CSE 486/586 Distributed Systems Security --- 2

Steve Ko  
Computer Sciences and Engineering  
University at Buffalo

CSE 486/586

### Recap

- Three types of functions
  - Cryptographic hash, symmetric key crypto, asymmetric key crypto
- Cryptographic hash
  - Easy to compute  $h(m)$
  - Hard to find an  $m$ , given  $h(m)$
  - Hard to find two values that hash to the same  $h(m)$
- How to find collisions?
  - Birthday paradox: for 50% prob. &  $m$  bits,  $\sim 2^{m/2}$  numbers
- Symmetric key crypto
  - MAC: Compute  $H = \text{AES}_k(\text{SHA1}(M))$  & Send  $\langle M, H \rangle$
- Asymmetric key crypto
  - Guarantees rely on computational hardness

CSE 486/586

2

### Recap

- MAC
  - Symmetric crypto
  - Verifies the authenticity of a message
  - Sender: compute  $H = \text{AES}_k(\text{SHA1}(M))$  & send  $\langle M, H \rangle$
  - Receiver: compute  $H' = \text{AES}_k(\text{SHA1}(M))$  & check  $H' == H$
- Digital Signatures
  - Asymmetric crypto
  - Signer: compute  $H = \text{RSA}_k(\text{SHA1}(M))$  & send  $\langle M, H \rangle$
  - Verifier: compute  $H' = \text{RSA}_k(H)$  & verify  $H' == \text{SHA1}(M)$
  - Not just integrity, but also authenticity

CSE 486/586

3

### Heard of Firesheep?

- Firesheep
  - A Firefox extension
  - A packet sniffer to intercept unencrypted cookies from certain websites (such as Facebook and Twitter)
  - Allows the user to take on the log-in credentials of the victim
- Solution?
  - Encrypt your traffic!
  - This is before facebook started using https, but now facebook uses https.

CSE 486/586

4

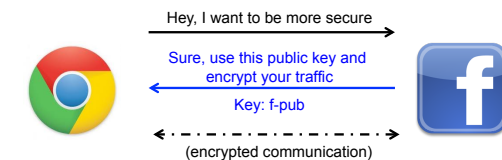
### “Securing” HTTP

- Threat model
  - Eavesdropper listening on conversation (confidentiality)
  - Man-in-the-middle modifying content (integrity)
  - Adversary impersonating desired website (authentication, and confidentiality)
- Enter HTTP-S
  - HTTP sits on top of secure channels
  - All (HTTP) bytes written to secure channel are encrypted and authenticated

CSE 486/586

5

### Encrypted Communication



- What is wrong with this?
  - How do you know you're actually talking to facebook and f-pub belongs to facebook?

CSE 486/586

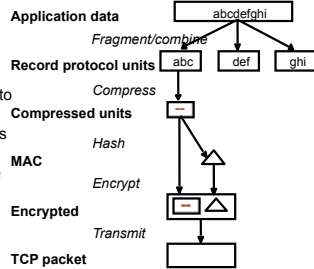
6



## TLS Record Protocol

- The record protocol takes an application message to be transmitted,

- fragments the data into manageable blocks,
- optionally compresses the data,
- computes a message authentication code (MAC),
- encrypts and
- adds a header.



CSE 486/586

13

## TLS Handshake Protocol

Cipher suite: a list of cryptographic algorithm supported by the client

### Phase 1: Establish security capabilities

ClientHello

ServerHello

Establish protocol version, session ID, cipher suite, compression method, exchange random values

### Phase 2: Server authentication and key exchange

Certificate

Certificate Request

ServerHelloDone

Optionally send server certificate and request client certificate

### Phase 3: Client authentication and key exchange

Certificate

Certificate Verify

Send client certificate response if requested

### Phase 4: Finish

Change Cipher Spec

Finished

Change Cipher Spec

Finished

Change cipher suite and finish handshake

The client sends a change Cipher Spec message and copies the pending CipherSpec into the current CipherSpec.

CSE 486/586

14

## CSE 486/586 Administrivia

- PA4 due Friday next week
- Final: 5/15 (Friday), 11:45am – 2:45pm
  - NSC 201

CSE 486/586

15

## Authentication

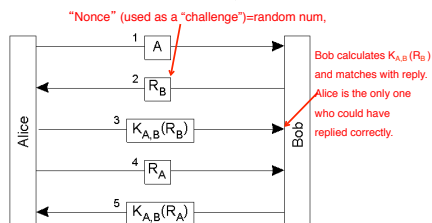
- Use of cryptography to have two principals verify each others' identities.
  - Direct authentication:** the server uses a shared secret key to authenticate the client.
  - Indirect authentication:** a trusted authentication server (third party) authenticates the client.
  - The authentication server knows keys of principals and generates temporary shared key (ticket) to an authenticated client. The ticket is used for messages in this session.
    - E.g., Verisign servers

CSE 486/586

16

## Direct Authentication

- Authentication with a secret key

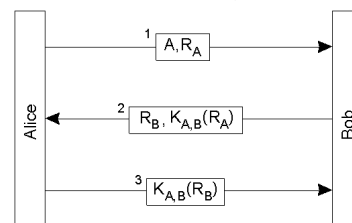


CSE 486/586

17

## “Optimized” Direct Authentication

- Authentication with a secret key with three messages

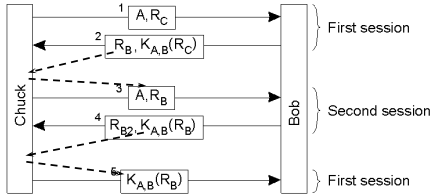


- Anything wrong with this?

CSE 486/586

18

## Reflection Attack



CSE 486/586

19

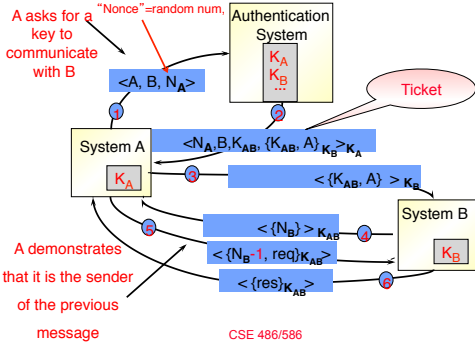
## Needham-Schroeder Authentication

- An **authentication server** provides secret keys.
  - Every client shares a secret key with the server to encrypt their channels.
- If a client A wants to communicate with another client B,
  - The server sends a key to the client A in **two forms**.
    - First, in a **plain form**, so that the client A can use it to encrypt its channel to the client B.
    - Second, in an **encrypted form** (with the client B's secret key), so that the client B can know that the key is valid.
  - The client A sends this encrypted key to the client B as well.
- Basis for Kerberos

CSE 486/586

20

## Needham-Schroeder Authentication



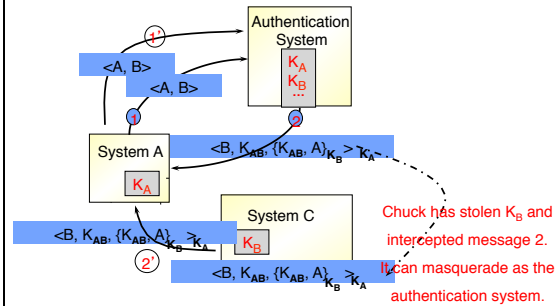
message

CSE 486/586

21

## Nonce $N_A$ in Message 1

Because we need to relate message 2 to message 1



CSE 486/586

22

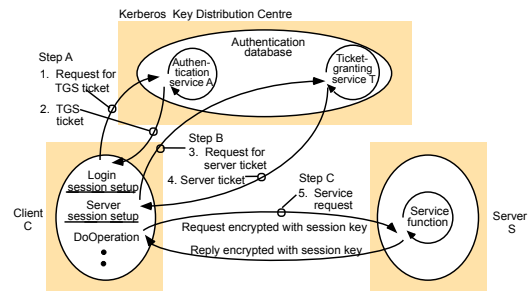
## Kerberos

- Follows Needham-Schroeder closely
- Time values used for nonces
  - To prevent replay attacks
  - To enforce a lifetime for each ticket
- Very popular
  - An Internet standard
  - Default in MS Windows

CSE 486/586

23

## Kerberos



CSE 486/586

24

## Summary

- Digital certificates
  - Binds a public key to its owner
  - Establishes a chain of trust
- TLS
  - Provides an application-transparent way of secure communication
  - Uses digital certificates to verify the origin identity
- Authentication
  - Needham-Schroeder & Kerberos

CSE 486/586

25

## Acknowledgements

- These slides contain material developed and copyrighted by Indranil Gupta (UIUC), Jennifer Rexford (Princeton) and Michael Freedman (Princeton).

CSE 486/586

26