

## Solution of Assignment #2, CS/191

Fall, 2014

1. Page 35, problem 10,

(0 points) (b)

$p$	$q$	$r$	$p \rightarrow q$	$q \rightarrow r$	$(p \rightarrow q) \wedge (q \rightarrow r)$	$p \rightarrow r$	$[(p \rightarrow q) \wedge (q \rightarrow r)] \rightarrow (p \rightarrow r)$
T	T	T	T	T	T	T	T
T	T	F	T	F	F	F	T
T	F	T	F	T	F	T	T
T	F	F	F	T	F	F	T
F	T	T	T	T	T	T	T
F	T	F	T	F	F	T	T
F	F	T	T	T	T	T	T
F	F	F	T	T	T	T	T

Since  $[(p \rightarrow q) \wedge (q \rightarrow r)] \rightarrow (p \rightarrow r)$  is always T, it is a tautology.

(0 points) (c)

sol:

$p$	$q$	$p \rightarrow q$	$p \wedge (p \rightarrow q)$	$[p \wedge (p \rightarrow q)] \rightarrow q$
T	T	T	T	T
T	F	F	F	T
F	T	T	F	T
F	F	T	F	T

Since  $[p \wedge (p \rightarrow q)] \rightarrow q$  is always T, it is a tautology.

(0 points) (d)

$p$	$q$	$r$	$p \vee q$	$p \rightarrow r$	$q \rightarrow r$	$(p \vee q) \wedge (p \rightarrow r) \wedge (q \rightarrow r)$	$[(p \vee q) \wedge (p \rightarrow r) \wedge (q \rightarrow r)] \rightarrow r$
T	T	T	T	T	T	T	T
T	T	F	T	F	F	F	T
T	F	T	T	T	T	T	T
T	F	F	T	F	T	F	T
F	T	T	T	T	T	T	T
F	T	F	T	T	F	F	T
F	F	T	F	T	T	F	T
F	F	F	F	T	T	F	T

Since  $[(p \vee q) \wedge (p \rightarrow r) \wedge (q \rightarrow r)] \rightarrow r$  is always T, it is a tautology.

2. (0 points), Page 35, problem 14.

sol:

$p$	$q$	$p \rightarrow q$	$\neg p \wedge (p \rightarrow q)$	$[\neg p \wedge (p \rightarrow q)] \rightarrow \neg q$
T	T	T	F	T
T	F	F	F	T
F	T	T	T	F
F	F	T	T	T

So  $[\neg p \wedge (p \rightarrow q)] \rightarrow \neg q$  is not a tautology.

3. (0 points), page 35, problem 18.

$$\begin{aligned}
 & p \rightarrow q \\
 \equiv & \neg p \vee q && \text{by the implication law (the first law in Table 7.)} \\
 \equiv & q \vee (\neg p) && \text{by commutative laws} \\
 \equiv & \neg(\neg q) \vee (\neg p) && \text{by double negation law} \\
 \equiv & \neg q \rightarrow \neg p && \text{by implication law}
 \end{aligned}$$

4. (0 points), page 35, problem 23.

$$\begin{aligned}
 & (p \vee q) \rightarrow r \\
 \equiv & \neg(p \vee q) \vee r && \text{by implication law} \\
 \equiv & (\neg p \wedge \neg q) \vee r && \text{by de Morgan's laws} \\
 \equiv & (\neg p \vee r) \wedge (\neg q \vee r) && \text{by distributive laws} \\
 \equiv & (p \rightarrow r) \wedge (q \rightarrow r) && \text{by implication laws twice}
 \end{aligned}$$

5. (0 points), page 35, problem 24.

$$\begin{aligned}
 & (p \rightarrow q) \vee (p \rightarrow r) \\
 \equiv & (\neg p \vee q) \vee (\neg p \vee r) && \text{by implication law, twice} \\
 \equiv & (\neg p \vee \neg p) \vee (q \vee r) && \text{by commutative and associative laws} \\
 \equiv & \neg p \vee (q \vee r) && \text{by Idempotent laws} \\
 \equiv & p \rightarrow (q \vee r) && \text{by implication law}
 \end{aligned}$$

6. (0 points), page 35, problem 26.

$$\begin{aligned}
 & \neg p \rightarrow (q \rightarrow r) \\
 \equiv & p \vee (q \rightarrow r) && \text{by implication law} \\
 \equiv & p \vee (\neg q \vee r) && \text{by implication law} \\
 \equiv & \neg q \vee (p \vee r) && \text{by commutative and associative laws} \\
 \equiv & q \rightarrow (p \vee r) && \text{by implication law}
 \end{aligned}$$

7. (0 points), page 35, problem 30. Show that  $(p \vee q) \wedge (\neg p \vee r) \rightarrow (q \vee r)$  is a tautology. sol:

$$\begin{aligned}
 & (p \vee q) \wedge (\neg p \vee r) \rightarrow (q \vee r) \\
 \equiv & \neg[(p \vee q) \wedge (\neg p \vee r)] \vee (q \vee r) && \text{by implication law} \\
 \equiv & [\neg(p \vee q) \vee \neg(\neg p \vee r)] \vee (q \vee r) && \text{by de Morgan's law} \\
 \equiv & [(\neg p \wedge \neg q) \vee (p \wedge \neg r)] \vee (q \vee r) && \text{by de Morgan's law} \\
 \equiv & [(\neg p \wedge \neg q) \vee q] \vee [(p \wedge \neg r) \vee r] && \text{by commutative and associative laws} \\
 \equiv & [(\neg p \vee q) \wedge (\neg q \vee q)] \vee [(p \vee r) \wedge (\neg r \vee r)] && \text{by distributive laws} \\
 \equiv & (\neg p \vee q) \vee (p \vee r) && \text{by negation and identity laws} \\
 \equiv & (\neg p \vee p) \vee (q \vee r) && \text{by communicative and associative laws} \\
 \equiv & T && \text{by negation and domination laws}
 \end{aligned}$$

8. (0 points), page 64, problem 6.

(d) sol: There is a student in your school who is enrolled in Math 222 and in CS 252.

(e) sol: There are two different students  $x$  and  $y$  such that if the student  $x$  takes the class  $z$ , then the student  $y$  also takes the class  $z$  for every classes  $z$ .

(f) sol: There are two different students  $x$  and  $y$  such that  $x$  takes  $z$  if and only if  $y$  takes  $z$  for all classes  $z$ .

9. page 65, problem 10,

(0 points)

(a) sol:  $\forall xF(x, Fred)$

(b) sol:  $\forall yF(Evelyn, y)$

(c) sol:  $\forall x\exists yF(x, y)$

(e) sol:  $\forall x\exists yF(y, x)$

(0 points)

(f) sol:  $\neg\exists x(F(x, Fred) \wedge F(x, Jerry))$

(i) sol:  $\neg\exists xF(x, x)$

10. page 67, problem 28,

(0 points)

(c) sol: True. Let  $x$  be 0.

(e) sol: True. Let  $y = 1/x$

(f) sol: False

(0 points)

(g) sol: True. Let  $y$  be  $1 - x$ .

(i) sol: False. Let  $x = 0$ . Then  $x + y = 2$  and  $2x - y = 1$  imply  $y = 2$  and  $y = -1$ . So for  $x = 0$  there exists no  $y$  such that  $x + y = 2$  and  $2x - y = 1$ .

(j) sol: True. Let  $z = (x + y)/2$

11. page 67, problem 30,

(0 points)

(b) sol:  $\exists x\forall y\neg P(x, y)$

(c) sol:  $\forall y(\neg Q(y) \vee \exists xR(x, y))$

(0 points)

(d) sol:  $\forall y(\forall x\neg R(x, y) \wedge \exists x\neg S(x, y))$

(e) sol:  $\forall y(\exists x\forall z\neg T(x, y, z) \wedge \forall x\exists z\neg U(x, y, z))$